Verifying Faust in Coq

Progress report

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CoqPL 2015



Music and PL?

Some Music	DSLs • DARMS			
 4CED Adagio Adal, AMPLE Arctic Arctic Arctic Arctic Arctic Arctic Arctic Canon CHANT Chack CLCE CMIX Chack CLCE CMIX Consolic CAMISIC Common Liqu Music Common Music Common Music 	 DCMP DMXX Elady Eatorpea Extempose Fases Fases FORMULA FORMULA FORMULA GROOVE GUIDO HARP HARP HARSI IMV 	 LPC Mars Masc Mat MideLogs MODE MOM MOM Most MUS10 MUS10 MUS50 MUS50 MUSCMD2ta MoseES MUSIC 10 MUSIC 10 MUSIC 11 MUSIC 160 	 MCL MUSIC III/IV/V MusicLago MusicLood MUSIC7 Musictex MUSIGOL Musictex MUSIGOL Musictex MIFF NOTELIST Nyquist OPAL OpenMusic OrganumI Outgetform Overtane FE 	 PLACODI PLAY3 PLAY3 PMX POCO POD5 POD7 PROD PROD Prodata PWGL Ravel SALIERI SCORE SCORE SCORFIE SCORFIE SCRIPT SIMDL SMOL SSP SSP ST
E Cound	 invokator KERN 	4BF # MUSIC 4F	# PILE	Supercollider Symbolic

Music and PL?

Some Music	DSLS * DARMS			
 # ACED # Adagiu # AddL # AddPLE # Arctic # Arctic # Arctichtang 	 DCMP DMX Elody ExAC Exterpro Enterpro Faux 	# LPC # Mars # MASC # Mas # MidLinp # MidLogs	 MCL MUSIC III/IV/V MusicLago Music1000 MUSIC7 	 PLACOM PLAY1 PLAY2 PMX POCO POD5 POD7
# Bang # Canon # CHANT # Obsock # CLCE # CMIX # Cmusic	Softw	vare ver	ification	toD redata • Royal • SALIERI • SCORE • ScoreFile
CMUSIC Common Lisp Music Common Music Common Music Common Music Notation Closend	GROOVE GUIDO HARP HARP Hastore HMSL INV invokator KENN	 MonetZata MusES MUSIC 10 MUSIC 11 MUSIC 340 MUSIC 48 MUSIC 48 MUSIC 47 	OPAL OpenMusic OrganumI Outperform Overtane PE Patcheork PILE	 SCRIPT SIREN SMDL SMOKE SSP SSP ST Supercollider Supercollider

Music and PL?

Some Music	DSLS . DARMS			
 4CED Adagio AddL AddPLE Aresic Autaklang 	 DCMP DMOX Elody EoAC Exterpro Exterpro Exterpro Fance 	# LPC # Mars # MASC # Mas # MidLisp # MidLogs	 MCL MUSIC III/IV/V MusicLago Music1000 MUSIC7 	 PLACOM PLAY1 PLAY2 PMX POCO POD6 POD7
Bang Canoo CHANT	Softv	vare ver	ification	ROD Heedata
Churck CLCE CMIX	* FORMES * FORMULA # Fague	Coq	P P P	 Ravel SALIERI SCORE ScoreFile
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= Csound	 invokator KERN 	4BF # MUSIC 4F	· PILE	# Supercollider # Symbolic

Faust

- ► Functional PL for digital signal processing.
- Synchronous paradigm, geared towards audio.
- Programs: circuits/block diagrams + feedbacks.
- Semantics: streams of samples.
- Efficiency is crucial.
- Created in 2000 by Yann Orlarey et al. at GRAME.
- Mature, compiles to more than 14 platforms.

Faust's Ecosystem

Users:

- ► Grame: Multiple projects, main developer.
- Stanford: Class/books on signal processing, STK instrument toolkit, Faust2android, Mephisto...
- Ircam: Acoustic libraries, effects libraries,...
- Guitarix, moForte guitar, etc...

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It has its market! Much easier than dwelling into C. Recent Events:

- Faust day at Stanford happened yesterday.
- ► Ongoing Faust program competition (€2,000 in prices).
- FEEVER project :)

Syntax and Well-Formedness

$$\begin{array}{ccc} \text{TERM} & \overline{\vdash !: 1 \rightarrow 0} & \text{ID} & \overline{\vdash _: 1 \rightarrow 1} \\ \\ \text{PAR} & \frac{\vdash f_1 : i_1 \rightarrow o_1 & \cdots & \vdash f_n : i_n \rightarrow o_n}{\vdash (f_1, \dots, f_n) : \sum_j^n i_j \rightarrow \sum_j^n o_j} \\ \\ \text{COMP} & \frac{\vdash f : i \rightarrow k & \vdash g : k \rightarrow o}{\vdash (f : g) : i \rightarrow o} \\ \\ \text{PAN} & \frac{\vdash f : i \rightarrow k & \vdash g : k \ast n \rightarrow o & 0 < k \land 0 < n}{\vdash f <: g : i \rightarrow o} \end{array}$$

Syntax and Typing

PL standard practice vs. what musicians want/imagine:



Figure 2: (B:C) sequential composition of *B* and *C*



Figure 3: sequential composition of *B* and *C* when k = 1

Feedbacks

$$\texttt{FEED} \; \frac{\vdash f: g_o + f_i \rightarrow g_i + f_o \quad \vdash g: g_i \rightarrow g_o}{\vdash f \sim g: f_i \rightarrow f_o}$$

Diagram for + \sim sin:



Synchronous semantics: execution in "ticks" + state.

Simple Low-pass Filter

```
smooth(c) = *(1-c) : + *(c);
process = smooth(0.9);
```



A More Real Example

Feedback Delay Networks:

fdnrev0(delays, BBS0, freqs, durs, loopgainmax, nonl) = (bus(2*N) :> bus(N) : delaylines(N)) (delayfilters(N,freqs,durs):feedbackmatrix(N)) with { delayval(i) = take(i+1,delays); delaylines(N) = par(i,N,(delay(dlmax(i),(delayval(i)-1)))); delayfilters(N,freqs,durs) = par(i,N,filter(i,freqs,durs)); feedbackmatrix(N) = bhadamard(N);vbutterfly(n) = bus(n) <: (bus(n):>bus(n/2)), ...) . . . };

A More Real Example









Does there exist any other programming language?

PHILOSOPHICAL

- Manual proofs starting to feel odd in PL.
- Motto: use Coq from the start.
- ► Goal: Try to develop in a reusable way.

PHILOSOPHICAL — MATHEMATICAL

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Current testing process of Faust programs: compare their output with MatLab's.

- Program correctness.
- Optimizations performed by the compiler are not well understood. Semantics trickier than it looks to the eye
- Explore the formalization of concepts from signal processing: Finite Impulse Response (FIR) filters, LTI theory, spectral analysis, Nyquist...

PHILOSOPHICAL — MATHEMATICAL PRACTICAL

Less effort than to build a custom analysis tool.

Applications:





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IMHO: **Robust Definitions and Standards** are crucial. Don't repeat the mistakes of the past

Some Properties of Interest

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In order to write the properties, we need a large support library

[bigops, intervals, trigonometry, Z-transforms, DTFT, ...]

Specifications of Filters

Difference equations:

$$y(n) = x(n) + x(n-1)$$

Impulse response:

$$H(z) = \frac{1 - z^{-2}}{1 - 2R\cos(\Theta_c)z^{-1} + R^2z^{-2}}$$

Two Poles Filter

$$H(z) = \frac{1 - z^{-2}}{1 - 2R\cos(\Theta_c)z^{-1} + R^2 z^{-2}}$$

with {

bw = 100; fr = 1000; g = 1; // parameters - see caption SR = fconstant(int fSamplingFreq, <math.h>); // Faust fn pi = 4*atan(1.0); // circumference over diameter R = exp(0-pi*bw/SR); // pole radius [0 required] A = 2*pi*fr/SR; // pole angle (radians) RR = R*R; firpart(x) = (x - x'') * g * ((1-RR)/2); feedback(v) = 0 + 2*R*cos(A)*v - RR*v';};

Finally! Let's Talk About Coq!

What we have built so far:

- Mathcomp allowed us to do a prototype in two weeks.
- New feedback reasoning rule & proof of soundness.
- Motivated by real use cases.
- Stateless logic & soundness (again, mc was key).
- Certified arity-checker, etc...

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Currently:

- Investigating more complex, time-aware logics.
- New semantics based on guarded recursion.

The Pieces of the Puzzle



The First Piece: Stream-based Semantics

- We ported [Boulmé, Hamon and Pouzet], with some problems with CoInductives.
- We switched to sequences, (similar to Auger's Lustre certified compiler).
- Didn't look into PACO/more advanced co-reasoning tools.

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Current approach: realizability semantics in guarded recursion style. Suggested simultaneously by A. Spiwak and A. Guatto [SYNCHRON 2014]:

$$\llbracket \vdash f : i \to o \rrbracket_W^n : \llbracket i \rrbracket^n \to \llbracket o \rrbracket^n$$

The Second Piece: Real Analysis

- Not in Mathcomp rcf good enough for experiments.
- Our typical case involves complex numbers, trigonometry and sums over infinite series.
- Think of proving Euler's formula:

$$e^{(j\Theta)} = \sin \Theta + j \cos \Theta$$

- Difficult to choose: Standard library? Coquelicot? C-CorN?
- Our feeling is that life is going to be very painful.

[We are ignoring floating point issues for now]

The Third Piece: Coq as a Tool

```
ile Edit Options Buffers Tools EasyCrypt Proof-General Help
                                                             Current goal (remaining: 2)
                                                             Type variables: <none>
                                                             &1 (left ) : VCGStep.vcg full
                                                             &2 (right) : VCGStep.vcg full s1
 equiv vcg step1: VCGStep.vcg full ~ VCGStep.vcg full s1
  true ==> =\{res\}.
                                                             pre = ={i} /\ i{1}
                                                             t =$ rmu
                                                                                                   (1) t =$ rmu
                                                             r =$ rmu
                                                                                                   (2) r =$ rmu
                                                             s1 =$ rmu
                                                                                                   (3) s1 =$ rmu
                                                             s2 =$ rmu
                                                                                                   (4) s2 =$ rmu
  + seq 1 1 : (i{1} /\ ={i,t}); first by auto.
                                                             insBr = i ? (t, r) : (r, t)
                                                                                                  (5) insBr = (t, r)
    seq 1 1 : (i{1} /\ ={i,t,r}); first by auto.
                                                             surrs = fst (vcg insBr (s1, s2) wt) (6) surrs = fst (vcg
    seg 1 1 : (i{1} /\ ={i.t.r.sl}): first by auto.
    seq 1 1 : (i{1} /\ ={i,t,r,s1,s2}); first by auto.
                                                             post =
     by wp; skip; progress; rewrite H.
                                                               (if i{1} then fst surrs{1} else snd surrs{1}) =
                                                               if i{2} then fst surrs{2} else snd surrs{2}
 + swap{1} 1 1.
    seq 1 1 : (! i{1} /\ ={i} /\ r{1} = t{2}); first by
auto.
     seq 1 1 : (! i{1} / ={i} / r{1} = t{2} / t{1} = re
42}); first by auto.
     seq 1 1 : (! i{1} /\ ={i,s1} /\ r{1} = t{2} /\ t{1}
s= r{2}) ; first by auto.
    seq 1 1 : (! i{1} /\ ={i,s1,s2} /\ r{1} = t{2} /\ t{
1] = r{2}) ; first by auto.
     by wp; skip; progress; rewrite H.
aed.
equiv vcg step2: VCGStep.vcg full s1 ~ VCGStep.vcg full •
                                                             U:%%- *goals*
                                                                                                  (EasyCrypt goals)
                                                              > Copyright (c) - 2012-2014 - IMDEA Software Institute a
s2 : true ==> ={res}.
                                                              >> Distributed under the terms of the CeCILL-C license
   swap{1} 5 2.
```

The Third Piece: Coq as a Tool

- Is building a verification tool on top of Coq feasible?
- We got some inspiration from domain-specific tools like EasyCrypt.
- Would our tool mature, we would certainly need to plug deeply into Coq's parsing/display routines.
- We still think this may be better than rewriting everything from scratch.
- Reduction woes make our life difficult.
- Automation: we will worry last.

Stability of Smooth

Recall the smooth program:

$$smooth(c) = *(1-c) : + *(c);$$

We want to prove stability, that is, bounded inputs produce bounded outputs, provided the coefficient c is in [0, 1]. We use the following logical rule (simplified):

$$\frac{ \left[\begin{array}{c} \models \psi(x_0) \\ \left\{ \gamma(i_1) \land \phi(i_2) \right\} & f \quad \{\psi(o)\} \\ \hline \left\{ \phi(i) \right\} & f \sim g \quad \{\psi(o)\} \\ \hline \left\{ \phi(i) \right\} & f \sim g \quad \{\psi(o)\} \end{array} \right\}$$

Stability of Smooth

Three VC in the proof:

```
by rewrite ?ler_wpmul2r ?ler_subr_addr ?add0r.
```

```
have Ha: a = a * c + a * (1 - c)
by rewrite -mulrDr addrC addrNK mulr1.
have Hb: b = b * c + b * (1 - c)
by rewrite -mulrDr addrC addrNK mulr1.
by rewrite Ha Hb !ler_add.
```

by rewrite ?ler_wpmul2r.

We pushed the VCs to Why3 with success. Interval technique ready to go into the main compiler.

Conclusions

- Young project, highly positive experience so far.
- First alpha "release" very near.
- Tons of related work, difficult to get a good perspective.
- Most challenging topic: real/complex analysis.
- Certified audio/dsp processing? (Do we need it?)
- All of the usual Coq caveats apply to us.
- What do *you* think?

Thanks!

Nyquist Theorem

Provided f_s is twice the highest frequency in V then:

$$V(t) = \sum_{n=-\infty}^{\infty} V[n] \frac{\sin[\pi f_s(t-nT_s)]}{\pi f_s(t-nT_s)}$$

where

 $f_s = 1/T_s$ sampling frequency V(t) value of signal at Time t $V[n] = V(nT_s)$ value of signal at Time $t = nT_s$